

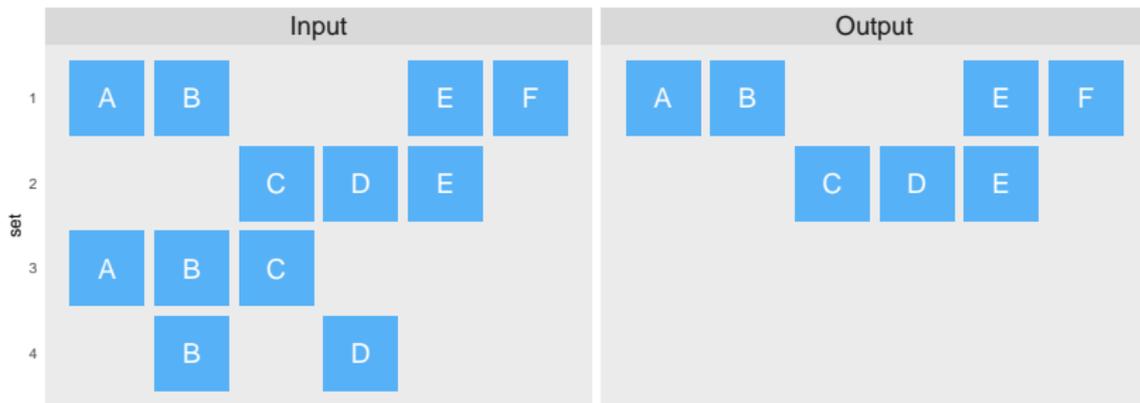
# Set Cover Problem

Input:  $\mathcal{S}$ , collection of sets  $S_1, \dots, S_n$  covering  $\mathcal{U}$ :

$$S_1 \cup S_2 \cup \dots \cup S_n = \mathcal{U}.$$

Output: Smallest subcollection from  $\mathcal{S}$ , covering  $\mathcal{U}$ .

# Set Cover: A Toy Example



# Set cover

- Fundamental problem in approximation algorithms with wide ranging applications; e.g. location planning, shift-planning and virus detection.
- Our application: Minimize number of hospitals, given every person in Germany can reach one hospital within 30 minutes.

# RcppGreedySetCover

- Optimal solution available via linear programming, not feasible for large problems.
- Alternative: Greedy approximation.
- No fast solution in R available → RcppGreedySetCover!
  - Fast due to use of `data.table` and Rcpp.

# Greedy Approximation: Algorithm

- Input:  $\mathcal{S} = \{S_1, \dots, S_n\}$ .
- Initialize  $\mathcal{C} \leftarrow \{\}$ ,  $\mathcal{T} \leftarrow \mathcal{S}$ .
- Repeat the following steps until  $\mathcal{C}$  is a cover of  $\mathcal{S}$ :
  1. Find the largest set of *uncovered* elements, say  $\Delta$ .
  2.  $\mathcal{C} \leftarrow \mathcal{C} \cup \Delta$ .
  3.  $\mathcal{T} \leftarrow \{T_1 \setminus \Delta, \dots, T_n \setminus \Delta\}$ .

# Greedy Approximation: Properties

- Tradeoff: Approximation error for speed / feasibility.
  - Error is bounded, approximation ratio depends on problem size.
- Vazirani 2001, p. 17: “[. . .], for the minimum set cover problem the obvious algorithm given above is essentially the best one can hope for.”

# Implementation

- Main requirement for containers: Efficient lookup and resizing.
  - No satisfactory solution in R available.
- Elements and sets are associated with integers  $0, 1, \dots$ 
  - Stored in `std::unordered_set<int>`. Grows / shrinks quickly.

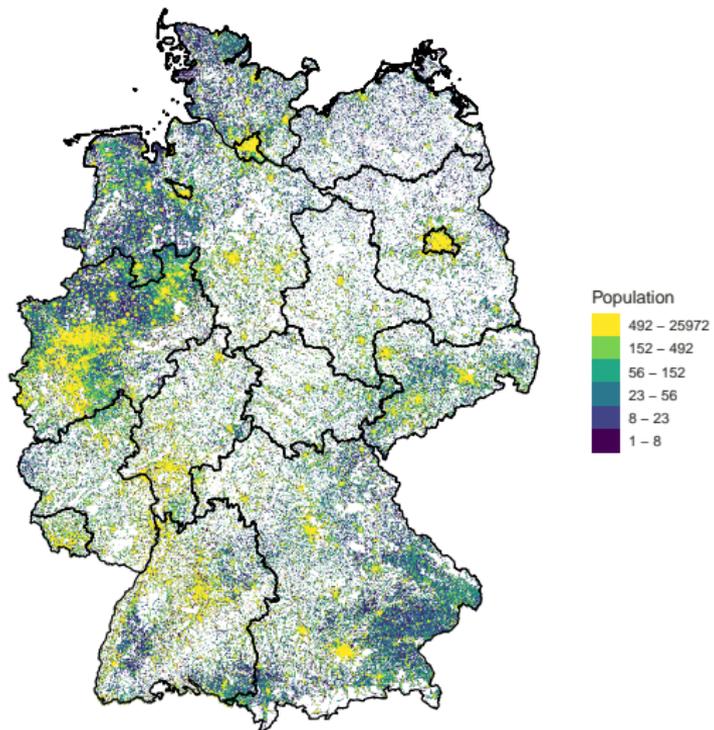
## Implementation: Objects

- Map sets to elements via `std::vector<.>`.
  - Integer-representation of sets for indexing.
  - $O(1)$  cost for element access.
- Map elements to sets via `std::unordered_map<int, .>`.
  - $O(1)$  average cost for access and removal.

# Implementation: Objects

- `boost::multi_index::multi_index_container` for set sizes.
- Custom interface:
  - Fast lookup of largest set size *and*
  - fast adjustment of set sizes.

# Application: Data



## Application: Data

Drivetimes for all populated grids-cells of 1km<sup>2</sup> raster in Germany within 40km radius, excluding drivetimes > 30 minutes.

```
head(D)
```

```
##           idm0           idm1 drivetime N_0 N_1
## 1: 4031_3109 4031_3110           157.2  92  23
## 2: 4031_3109 4031_3111           341.1  92  50
## 3: 4031_3109 4032_3108           198.8  92 166
## 4: 4031_3109 4032_3109           125.0  92 258
## 5: 4031_3109 4032_3111           298.7  92 244
## 6: 4031_3109 4033_3104           870.2  92 126
```

```
nrow(D)
```

```
## [1] 164114074
```

# Application

- Input:  $N \times 2$  tidy data.frame. Sets are in the first, elements in the second column.

```
library(RcppGreedySetCover)
system.time(
  Res <- greedySetCover(D[,c("idm0", "idm1")])
)
```

```
## 100% covered by 867 sets.
```

```
##      user  system elapsed
## 298.57    8.89   307.57
```

# Application

- Output analogous to input.

```
head(Res)
```

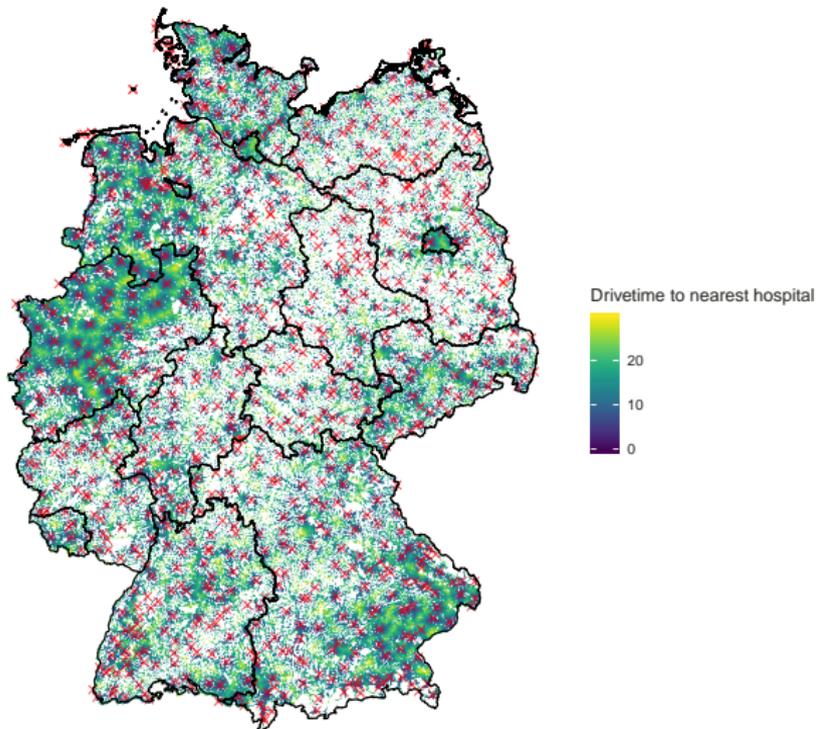
```
##           idm0           idm1
## 1: 4041_3197 4041_3189
## 2: 4041_3197 4041_3190
## 3: 4041_3197 4042_3189
## 4: 4041_3197 4046_3199
## 5: 4041_3197 4052_3180
## 6: 4046_3075 4040_3086
```

```
# Sanity check, TRUE if solution is a cover:
setequal(Res$idm1,D$idm1)
```

```
## [1] TRUE
```

# Application: Result

- Hospital marked by red X.



## Future improvements

1. Speed up implementation.
2. Get rid of `data.table` dependency.
3. Extend to weighted set cover / capacitated set cover.